Exceptions

Up to this point in the semester, we have assumed that the input to our functions are always correct, and have thus not done any error handling. However, functions can often have large domains, and we want our functions to handle erroneous input gracefully. This is where exceptions come in.

Exceptions provide a general mechanism for adding error-handling logic to programs. Raising an exception is a technique for interrupting the normal flow of execution in a program, signaling that some exceptional circumstance has arisen.

An **exception** is an object instance of a class that inherits, either directly or indirectly, from the BaseException class. The following is an example of how to raise an exception:

```
>>> raise Exception('An error occurred')
Traceback (most recent call last):
File "<stdin>", line 1, in <module>
Exception: An error occurred
```

Notice how the string 'An error occurred' is an argument to the Exception object being created, and the string is part of what Python prints out in response to the exception being raised. If the exception is raised while within a try statement, then the interpreter will immediately look for an except statement that handles the type of exception being raised. try and except statements allow programs to respond to unexpected arguments and other errors gracefully, rather than terminating entirely.

Here's is how to structure try and except statements:

Questions:

1. Fill in all the blanks to produce the desired output:

2. Write the function <code>safe_square</code> that uses exceptions to print 'Incorrect argument type' when anything other than an instance of the <code>int</code> or <code>float</code> class is given as an argument. Otherwise, <code>safe_square</code> should multiply the argument by itself. A useful fact is that a <code>TypeError</code> is raised when * is given incorrect arguments.

3. Predict the output of each of the following lines, assuming safe square is implemented as described in question 2.

```
>>> safe_square('hello')
Incorrect argument type
>>> safe_square('hello' * 5)
Incorrect argument type
>>> safe_square('hello' * 'hello')
TypeError: can't multiply sequence by non-int of type 'str'
>>> safe_square(1 * 2.5)
6.25
>>> safe_square(1 / 0)
ZeroDivisionError: division by zero

4. Consider the following:
>>> def hello_append(lst):
    try:
        if type(lst) != list:
```

raise HelloAppendError(lst)

print('type of arg was', type(e.contents))

Create a HelloAppendError class that would create the following output given the definition of hello_append above:

lst.append('hello')
except HelloAppendError as e:

When writing big programs, it's usually a good idea to have robust error handling logic that details the state of the program so that you, as the programmer, can debug your code more effectively. Additionally, by gracefully responding

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to errors, your program could potentially continue to run even when errors or unexpected input occur. In this class, we are primarily interested in using exceptions as part of the error handling logic that we will implement when we start creating interpreters in Python.

Orders of Growth

(Notes developed based on those of George Wang, Jon Kotker, Seshadri Mahalingam, Chung Wu, and Justin Chen)

When we talk about the efficiency of a procedure (at least for now), we're often interested in how much more expensive it is to run the procedure with a larger input. That is, as the size of the input grows, how do the speed of the procedure and the space its process occupies grow?

For expressing all of these, we use what is called the Big-Theta notation. For example, if we say the running time of a procedure foo is in $\Theta(n^2)$, we mean that the time it takes to process the input grows as the square of the size of the input. More generally, we can say that foo is in some $\Theta(f(n))$ if there exist some constants k1 and k2 such that and some constants c1 and c2 such that,

k1 * f(n) < running time of foo for some n > c1, and <math>k2 * f(n) > running time of foo for some n > c2

To prove, then, that foo is in $\Theta(f(n))$, we only need to find constants k1 and k2 where the above holds. Fortunately for you, in 61A, we're not that concerned with rigor, and you probably won't need to know exactly how to do this (you will get the painful details in 61B!) What we want you to have in 61A, then, is the intuition of guessing the orders of growth for certain procedures.

Kinds of Growth

Here are some common ones: $\Theta(1)$ – constant time (takes the same amount of time regardless of input size); $\Theta(\log n)$ – logarithmic time; $\Theta(n)$ – linear time; $\Theta(n^2)$, $\Theta(n^3)$, etc – polynomial time; $\Theta(2^n)$ – exponential time ("intractable"; these are really, really horrible).

Orders of Growth in Time

"Time", for us, basically refers to the number of recursive calls or the number of times the suite of a while or for loop executes. Intuitively, the more recursive calls we make, the more time it takes to execute the function.

- If the function contains only primitive procedures like + or *, then it is constant time $-\Theta(1)$.
- If the function is recursive, you need to:
 - a. Count the number of recursive calls there will be given input n
 - b. Count how much time it takes to process the input per recursive call

The answer is usually the product of the above two. For example, given a fruit basket with 10 apples, how long does it take for me to process the whole basket? Well, I'll recursively call my eat procedure which eats one apple at a time (so I'll call the procedure 10 times). Each time I eat an apple, it takes me 30 minutes. So the total amount of time is just 30*10 = 300 minutes!

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- If the function contains calls of helper functions that are not constant-time, then you need to take the orders of growth of the helper functions into consideration as well. In general, how much time the helper function takes would be factored.
- When we talk about orders of growth, we don't really care about constant factors. So if you get something like $\Theta(1000000n)$, this is really $\Theta(n)$. We can also usually ignore lower-order terms. For example, if we get something like $\Theta(n^3 + n^2 + 4n + 399)$, we take it to be $\Theta(n^3)$.

Questions:

Give the order of growth in time for the following functions:

```
1. def sum of factorial(n):
      if n == 0:
            return 1
      else:
            return factorial(n) + sum of factorial(n-1)
\Theta(n^2)
2. def fib iter(n):
      prev, curr, i = 0, 1, 1
      while i < n:
            prev, curr = curr, prev + curr
      return curr
0 (n)
3. \text{ def mod } 7(n):
      if n % 7 == 0:
            return 0
      else:
           return 1 + mod_7(n-1)
\Theta(1)
4. Given:
def bar(n):
      if n % 2 == 1:
            return n + 1
      return n
def foo(n):
      if n < 1:
            return 2
      if n % 2 == 0:
            return foo (n-1) + foo (n-2)
      else:
            return 1 + foo(n-2)
```

What is the order of growth of foo (bar (n))?

 $\Theta(n^2)$

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Calculator

We are beginning to dive into the realm of *interpreting* computer programs. In order to do so, we'll have to examine some new programming languages. The *Calculator* language, invented for this class, is the first of these examples.

Our Calculator language is actually nearly identical to Python, except it only handles arithmetic operations - add, sub, mul, div, etc. Also unlike Python's arithmetic operators, we'll let Calculator's operators each take an arbitrary number of arguments.

Here's a few examples of Calculator in action:

```
calc> 6
6
calc>mul()
1
calc>add(1, mul(3, sub(3, 7)))
-11
```

Our goal now is to write an interpreter for the Calculator language. The job of an interpreter is, given an expression, evaluate its meaning. So let's talk about expressions.

Representing Expressions

When we type a line at the Calculator prompt and hit enter, we've just sent an expression to the interpreter. We can represent an Expression as an object:

```
class Exp(Object):
    def __init__(self, operator, operands):
        self.operator = operator
        self,operands = operands

def __repr__(self):
        return 'Exp({0}, {1})'.format(repr(self.operator), repr(self.operands)))

def __str__(self):
        operand_strs = ', '.join(map(str, self.operands))
        return '{0}({1})'.format(self.operand, operand strs)
```

Don't worry about what repr and str do for an Exp right now, the most important thing to note is that every Exp by its operator and operands. Also another thing to note: an Exp's operands are always themselves Exps as well.

Example: If I wanted to represent the Calculator expression add (2, 3), I would do this by calling the Exp constructor as follows: Exp ('add, [2, 3]).

Our Calculator language is only concerned with two kinds of expressions: numbers, which are *self-evaluating* expressions, and call expressions, which involve an operator acting on some number of arguments, each of which are expressions.

What does it mean for a number to be self-evaluating? If we evaluate an expression that is a number, the value of the expression (i.e. the result returned by evaluating it) is that number. Simple!

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What about evaluating call expressions? We can follow this straightforward two-step process.

- 1. Evaluate the operands.
- 2. Apply the operator, to the arguments (the values of the operands) (Sound familiar?)

Using this two-step process, we can interpret any Calculator expression. The first step will be handled by a function called calc eval, and the second step will be handled by a function called calc apply.

```
Here's calc_eval:

def calc_eval(exp):
    if type(exp) in (int, float): # if the expression is a number
        return exp
    else:
        arguments = list(map(calc_eval, exp.operands)) #evaluate the operands
        return calc apply(exp.operator, arguments) #apply the function to operands
```

As you can see, all we've done is follow the rules of evaluation outlined above. If the expression is a number, return it. Else, evaluate the operands and apply the operator the evaluated operands.

How do we apply the operator? With calc apply:

```
def calc_apply(operator, args):
    if operator in ('add', '+'):
        return sum(args)
    if operator in ('sub', '-'):
        if len(args) == 0:
            raise TypeError(operator + ' requires at least 1 argument')
        if len(args) == 1:
            return -args[0]
        return sum(args[0], [-args for args in args[1:]])
    if operator in ('mul', '*'):
        return reduce(mul, args, 1)
```

Depending on what the operator is, we can match it to a corresponding Python call. Each conditional in the function above corresponds to the application of one operator.

Something very important but may not have been obvious: calc_exp deals with *expressions*, calc_apply deals with *values*.

Questions

Let's say we want to make the variable a contain the object representation of the Calculator expression: add (4, 5, mul (3, 2)). Fill in the blank:

```
>>>a = Exp('add', [4, 5, Exp('mul', [3, 2])])
```

Suppose we typed the following expression into the calc interpreter:

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```
calc> add(1, mul(3, sub(3, 7)))
```

How many calls to calc_eval does this generate? How many calls to calc_apply?

7 calls to calc eval, 3 calls to calc apply.

The above implementation of the calc_exp and calc_apply does not handle calls to div. Add to them so we can evaluate Calculator div calls:

Example:

```
calc>div(8, 4)
2
...
if operator in ('div', '/'):
    if len(args) != 2:
        raise TypeError(operator + 'requires 2 arguments')
    num, denom = args
    return num / denom
```